

Bidder Valuations

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Introduction

Auction formats

- In first price sealed bid (FPSB) auctions the highest bidder wins and pays his bid.
- In second price sealed bid auctions (SPSB) the highest bidder wins and pays the bid of highest losing bidder.
- In Dutch auctions (reducing the price until a player accepts the offer) only the winning bid is ever observed; Dutch auctions are strategically equivalent to FPSB auctions.
- In Japanese (button) auctions players exit as the auctioneer raises the price and the winner pays the price at which the only other remaining bidder exits.
- Note that players update their information sets in Japanese auctions so are not necessarily strategically equivalent to SPSB auctions.

Example

Independent and identically distributed private values in a first price sealed bid auction

- We first consider a first price sealed bid (FPSB) auction for N players with independent private values (IPV).
- By FPSB we mean that each player $n \in \{1, \dots, N\}$ simultaneously submits a bid denoted by $b_n \in \mathbf{R}^+$, and that the player submitting the highest bidder is awarded the (single) object up for auction, and pays what he or she bid.
- By IPV we mean that for each $n \in \{1, \dots, N\}$ the value of owning the object is v_n where $v_n \in \mathbf{V}$ independently drawn from a common distribution, $F(v)$.

Example

Best replies in equilibrium

- Let $W(b)$ denote the probability of winning the auction with bid b . That is:

$$W(b) \equiv \Pr \{b_k \leq b \text{ for all } k = 1, \dots, N\}$$

- Then the maximization problem faced by player n can be written as:

$$\max_b (v_n - b) W(b)$$

- The first order condition (FOC) is:

$$(v_n - b_n) W'(b_n) - W(b_n) = 0 \quad (1)$$

- The second order condition (SOC) of the optimization problem is:

$$\begin{aligned} \frac{\partial}{\partial b} FOC &\equiv \frac{\partial}{\partial b} [(v - b) W'(b) - W(b)] \\ &= (v - b) W''(b) - 2W'(b) \\ &\equiv SOC \\ &< 0 \end{aligned}$$

Example

Pure strategy best replies are increasing in valuations

- Totally differentiating the FOC with respect to b and v yields:

$$0 = W'(b_n) dv_n + [(v_n - b_n)W''(b_n) - 2W'(b_n)] db_n$$

and hence:

$$\frac{db_n}{dv_n} = \frac{-W'(b_n)}{(v_n - b_n)W''(b_n) - 2W'(b_n)} > 0$$

because $W'(b_n) > 0$ and the denominator of the quotient is the SOC.

- We infer that if players are in a pure strategy equilibrium with an interior solution, then b_n is increasing in v_n .

Example

Bayesian Nash Equilibrium with monotone bidding

- From now on we assume that players are in a (pure strategy) Bayesian equilibrium with bids that are monotone increasing in valuations.
- That is we consider Bayesian Nash Equilibrium (BNE) in which bidders follow a strategy $\beta : \mathbf{V} \rightarrow \mathbf{B} \equiv [0, \infty)$ where $\beta(v)$ is increasing in v .
- Then $\beta(v)$ has an inverse, which we denote by $\alpha : \mathbf{B} \rightarrow \mathbf{V}$ such that $\alpha[\beta(v)] = v$ for all v .
- Letting $G(b)$ denote the distribution of bids, it follows that:

$$W(b) \equiv \Pr \{b_k \leq b_n \text{ for all } k = 1, \dots, N\} = G(b_n)^{N-1}$$

- From the monotonicity property of the BNE:

$$G(b) = F(\alpha(b))$$

Example

Identification when all bids are observed from the probability of winning

- Assume our data set consist of all the bids recorded in I auctions in which the same equilibrium is played.
- Let b_n^i for $n \in \{1, \dots, N\}$ and $i \in \{1, \dots, I\}$ denote the n^{th} bid in the i^{th} auction.
- The probability of winning the auction, $W(b)$, and its derivative $W'(b)$ are identified.
- We rewrite the FOC, Equation (1) as:

$$v_n^i = b_n^i + \frac{W(b_n^i)}{W'(b_n^i)} \quad (2)$$

- This shows v_n^i is identified, and therefore so is $F(v)$.

Example

A nonparametric estimator of the valuations

- Suppose there are a total of NI bids in the I auctions.
- Then v_n^i can be estimated nonparametrically with:

$$v_n^i = b_n^i + \widehat{W}(b_n^i) \left/ \frac{\partial \widehat{W}(b_n^i)}{\partial b} \right. \quad (3)$$

where:

$$\widehat{W}(b) = \frac{\sum_{m=1}^{M \equiv NI} I \{b_m \text{ is a winning bid}\} K\left(\frac{b-b_m}{h_M}\right)}{\sum_{m=1}^{M \equiv NI} K\left(\frac{b-b_m}{h_M}\right)}$$

and appealing to Collomb (1977; C.R. Acad. Sci. Paris 285, 289-292):

- $K(\cdot)$ is uniformly continuous, has bounded variation, and is integrable, with $K(-\infty) = K(\infty) = 0$ and $0 < K(0) < \infty$;
- for example $K(\cdot)$ is a standard normal pdf;
- $h_M \rightarrow 0$ and $M^{-1}h_M^{-1} \log M \rightarrow 0$ as $M \rightarrow \infty$.

Example

Identification when all bids are observed from the bidding distribution

- Alternatively note that the probability distribution of bids and its density, $G(b)$ and $G'(b)$, are identified.
- But the probability n wins with b_n is:

$$W(b_n) = G(b_n)^{N-1}$$

implying

$$W'(b_n) = (N-1) G(b_n)^{N-2} G'(b_n)$$

- We rewrite the FOC, Equation (1) as:

$$v_n^i = b_n^i + \frac{W(b_n^i)}{W'(b_n^i)} = b_n^i + \frac{G(b_n^i)}{(N-1) G'(b_n^i)} \quad (4)$$

- This shows v_n^i and hence $F(v)$ can also be directly identified off the bidding distribution $G(b)$.

Example

The distribution of winning bids

- Now suppose our data set consist of only the winning bid recorded in I auctions in which the same equilibrium is played.
- Let b^i for $i \in \{1, \dots, I\}$ denote the winning bid in the i^{th} auction.
- Thus the distribution of winning bids, denoted by $H(b^i)$, is identified.
- Since the winning bid is defined as the highest one, $H(b)$ is just the probability that all the bids are less than b , implying:

$$H(b) = \Pr \{ b_n^i \leq b \text{ for all } n = 1, \dots, N \} = G(b)^N$$

- Consequently:

$$G(b) = H(b)^{\frac{1}{N}} \quad (5)$$

and

$$G'(b) = \frac{1}{N} H(b)^{\frac{1}{N}-1} H'(b) \quad (6)$$

- This shows the bidding distribution is identified from the data generating process of the winner's bid.

Example

Identification when only the winning bid is observed

- Substituting Equations (5) and (6) back into Equation (4) gives:

$$v^i = b^i + \frac{G(b^i)}{(N-1)G'(b^i)} = b^i + \frac{NH(b)}{(N-1)H'(b)}$$

- This identifies the winning valuations, and hence their distribution, denoted by $F_W(v)$.
- But the distribution of the winning valuations is a one to one mapping of the distribution of all the valuations:

$$F_W(v) = \Pr \{v_n \leq v \text{ for all } n = 1, \dots, N\} = F(v)^N$$

- Therefore $F(v)$ is identified off the winning bids alone using the equation:

$$F(v) = F_W(v)^{\frac{1}{N}}$$

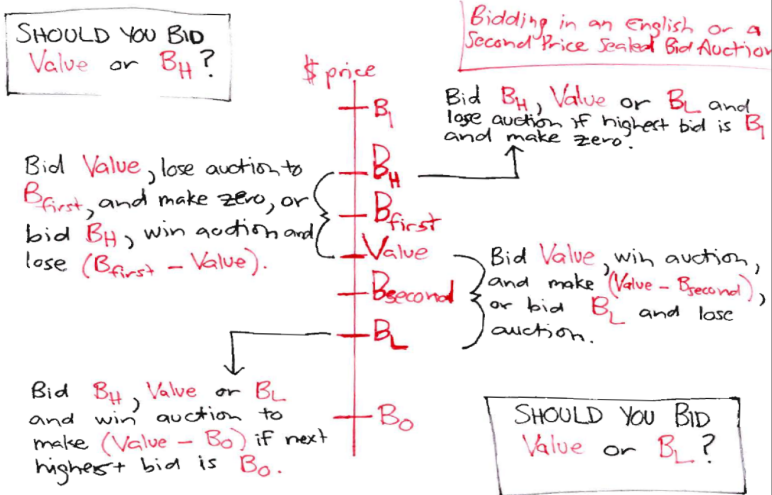
Another Example

A second price sealed bid (SPSB) auction with private values

- Now suppose as before:
 - each bidder knows her own valuation;
 - makes sealed bid (that is bids simultaneously).
- But instead of a FPSB auction, consider a SPSB auction, where the highest bidder wins the auction but only pays the second highest bid.
- Now it is a weakly dominant strategy for (each) n to bid her expected valuation, v_n .
- Intuitively, compared with bidding v_n :
 - bidding more implies winning some auctions that yield negative expected value, but leaves unchanged the expected value of any other auction that would be won;
 - bidding less implies losing some auctions that yield positive expected value, but leaves unchanged the expected value of any other auction that she would win.

Another Example

A picture proof



Another Example

Distribution of the second highest valuation

- Let $F(v)$ denote the distribution of valuations as before.
- Note first the obvious point that because players bid their valuations in SPSB auctions with private valuations, $F(v)$ is trivially identified if all the bids are observed.
- Now suppose only the winning price is observed.
- Then the probability distribution of the second highest valuation, which we now denote by $F_{N-1,N}(v)$, is identified.

Another Example

Identification when only the winning bid is observed

- More generally, let $F_{i,N}(v)$ denote the distribution of the i^{th} order statistic.
- Then it can be shown (see Arnold, Balakrishnan and Nagaraja, 1992 for example) that:

$$F_{i,N}(v) = \frac{N!}{(N-i)!(i-1)!} \int_{\underline{v}}^{F(v)} t^{i-1} (1-t)^{N-i} dt$$

- Note that \underline{v} is identified (and a consistent estimate is the lowest winning bid observed in the data).
- Also:

$$\frac{\partial F_{i,N}(v)}{\partial v} = \frac{N!}{(N-i)!(i-1)!} F(v)^{i-1} [1 - F(v)]^{N-i}$$

Theoretical Foundations

Notation and terminology for sealed bid auctions

- There are N players, risk neutral bidders.
- Bidder n :
 - has valuation v_n , the utility gain from winning the auction.
 - receives signal x_n , which without loss of generality we normalize, setting $x_n \equiv E[v_n | x_n]$. Let $x \equiv (x_1, \dots, x_N)$.
- We often assume $y \equiv (v_1, \dots, v_N, x)$ is affiliated, higher realizations of one component associated with higher realizations of the others.
- This means for random variable Y with pdf $f_Y(y)$, where \vee (\wedge) denotes the component wise maximum (minimum):

$$f_Y(y \vee y') f_Y(y \wedge y') \geq f_Y(y) f_Y(y')$$

- Bidders have:
 - private valuations if $E[v_n | x] = x_n$;
 - common valuations if $E[v_n | x]$ is strictly increasing in all its conditioning components (x_1, \dots, x_N) .
 - pure common values if $E[v_n | x] = E[v_{n'} | x]$ for all n and n' .

Theoretical Foundations

Equilibrium best responses in second price auctions with private values

- The literature focuses on perfect Bayesian equilibria in weakly undominated pure strategies (Athey and Haile, 2006).
- Let $b_n \equiv \beta_n(x_n, N)$ denote the equilibrium strategy of bidder n .
- In a second price auction with private values, it is a weakly dominant strategy for (each) n to bid his expected valuation, setting:

$$\beta_n(x_n, N) = x_n \equiv E[v_n | x_n]$$

- Note the same logic applies to n individually if $v_n = x_n$, regardless of the correlation structure of y and the other bidders' information.

Theoretical Foundations

Equilibrium best responses in first price auctions with private values

- In a private value FPSB auction denote the bid distribution function for the maximum equilibrium bid of the n^{th} bidder's rivals, conditional on the signal of n , as:

$$G_{m_n}(m | x_n, N) = \Pr \left[\max_{n' \in N \setminus n} \{b_{n'}\} \leq m \mid x_n, N \right]$$

- Then b_n solves:

$$b_n = \arg \max_b \int_{-\infty}^b (x_n - b) G'_{m_n}(m | x_n, N) dm$$

- The first order condition is:

$$x_n = b_n + \frac{G_{m_n}(b_n | x_n, N)}{G'_{m_n}(b_n | x_n, N)}$$

- Note this FOC reduces to (2) when $v_n = x_n$ and the valuations of the bidders are *iid*; in any case both $W(b_n)$ and $G_{m_n}(b_n | x_n, N)$ represent the probability of n winning the auction with bid b_n .

Theoretical Foundations

Equilibrium best responses in first price auctions with common values

- At a superficial level, this first order condition takes a similar form in a common value auction. Define:

$$v_n(x_n, x'_n, N) = E \left[v_n \left| x_n \text{ and } \max_{n' \in N \setminus n} \{b_{n'}\} = \beta_n(x'_n, N) \right. \right]$$

- Similar to the private values case b_n solves:

$$b_n = \arg \max_b \int_{-\infty}^b [v_n(x_n, \beta_n^{-1}(m, N), N) - b] G'_{m_n|b}(m | x_n, N) dm$$

- The first order condition is:

$$v_n(x_n, x_n, N) = b_n + \frac{G_{m_n|b}(b_n | x_n, N)}{G'_{m_n|b}(b_n | x_n, N)}$$

Identification in FPSB Auctions with Private Values

When all the bids are observed

- Assume $x_n = v_n$. From the first order condition:

$$x_n = b_n + \frac{G_{m_n}(b_n | x_n, N)}{G'_{m_n}(b_n | x_n, N)}$$

- Recall from its definition that $G_{m_n}(b_n | x_n, N)$ is the probability that n wins the auction with b_n :

$$G_{m_n}(b_n | x_n, N) = \Pr \left[\max_{n' \in N \setminus n} \{b_{n'}\} \leq b_n | x_n, N \right]$$

- Thus if all the bids are observed then $G_{m_n}(b_n | x_n, N)$ is identified.
- Hence v_n is identified (for all bidders in each sampled auction).
- Therefore the probability distribution of (v_1, \dots, v_N) in this specialization is identified for any correlation structure.

Identification Fails in Common Value FPSB Auctions

When all the bids are observed

- Recall that we defined:

$$v_n(x_n, x'_n, N) = E \left[v_n \mid x_n \text{ and } \max_{n' \in N \setminus n} \{b_{n'}\} = \beta_n(x'_n, N) \right]$$

and derived:

$$v_n(x_n, x_n, N) = b_n + \frac{G_{m_n}(b_n | x_n, N)}{G'_{m_n}(b_n | x_n, N)}$$

- The basic problem is that conditional on N the RHS gives a number for each n , but the LHS is not a primitive of the model.
- Note that every common value model is observationally to a private value model found by setting $v_n = v_n(x_n, x'_n, N)$.
- Thus two common value models with the same $v_n(x_n, x'_n, N)$ are (also) observationally equivalent.